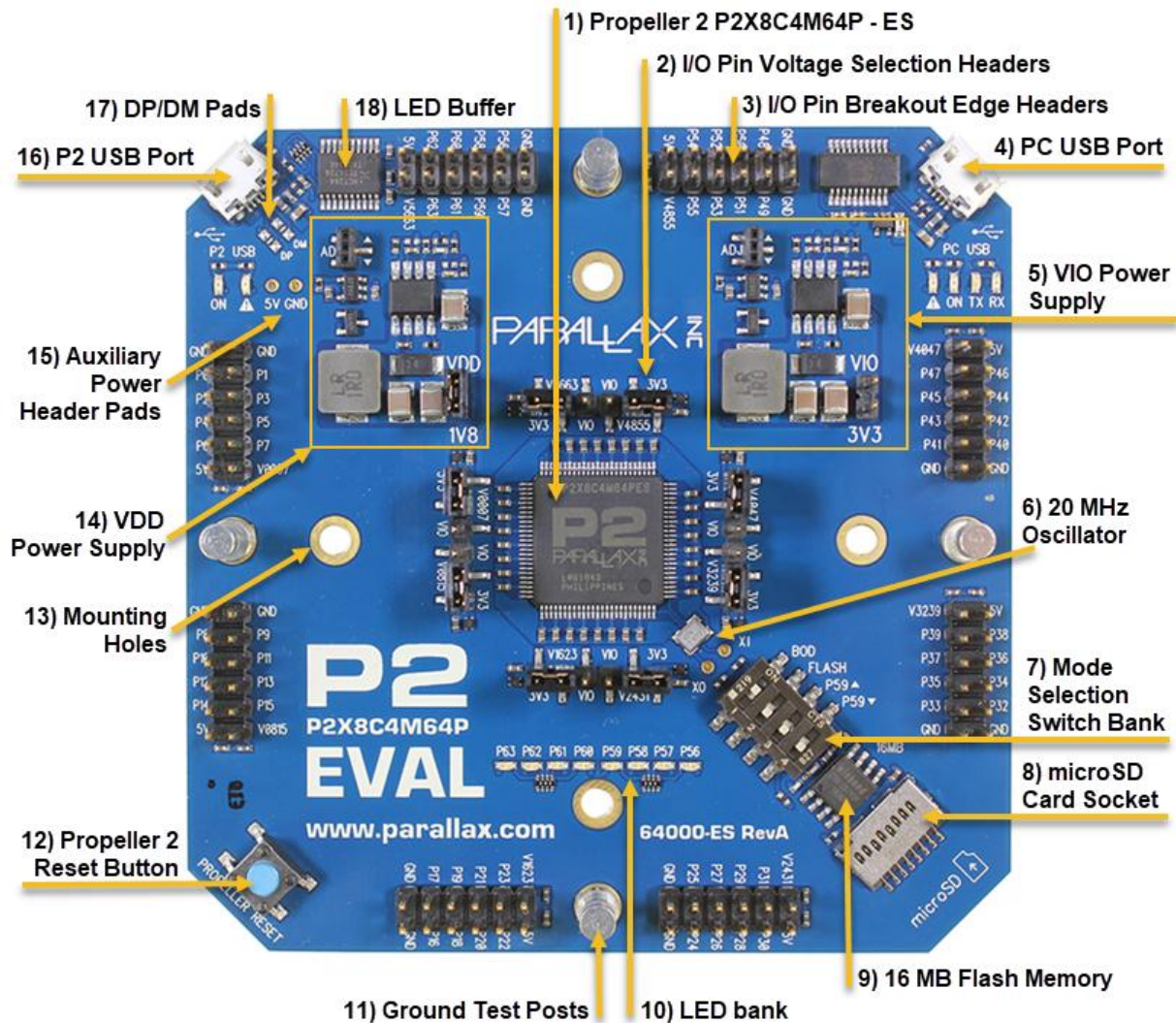


## Differences between P1 and P2 Pasm

This is a second tutorial regarding propeller pasm programming. I have had since December the P2 Engineering Sample. There are many differences between both the P1 and P2.

The P2 is 64 bit and 16 cores and 64 pins.



My first tutorial as well as with this one was for myself, but I shared it to help anybody who was confused and frustrated with getting into assembly language programming. I am adding to it some of the basic differences using the same approach.

I will use the same approach. Leaving the first tutorial in place, the P2 pasm code will be shown after page 22.

Still no Blinky lights in the beginning!!!! They will come later for both P1 and P2 heading towards running stepper motors.

## Differences between P1 and P2 Pasm

The first thing one will need is a copy of the propeller manual that is in the propeller tool and can be found here: <https://www.parallax.com/product/122-32000>.

Here is a link to Jeff Martin's webinar I uploaded to YouTube:  
<https://www.youtube.com/watch?v=OZHuWYW3o1A>

The first exercise will encompass passing variables from a spin method to a pasm method and back.

This the first piece of code that I came up with. There may be better ways to do this so bear with me.

I setup two global variables one for the spin method and the other for the pasm method. A five second waitcnt is used so as to have time to open the serial terminal when launching the code.

In order to launch the pasm code into a new cog this command is needed:

cognew(@asm,@datavar). The cognew means open the next cog, the @asm is the beginning of the assembly routine and the @datavar is the address of the first global variable.

```
1
2
3 {{ Tutorial 1 how to pass a number variable from spin to pasm and back, this works for numbers
4 from 0 to 256, bigger numbers in a later tutorial}}
5
6
7 CON
8     _clkmode = xtal1 + pll16x
9     _xinfreq = 6_250_000    'MY BOARD AT 100MHZ  DIFFERENT CRYSTAL
10    _xinfreq = 5_000_000    'QUICKSTART 80 MHZ  NORMAL CRYSTAL
11
12
13
14 obj
15
16     pst:"parallax serial terminal"
17
18 var
19     long datavar
20     long answwer
21
22
```

Figure 1

The next steps are to start the serial terminal wait five seconds to allow one to open the serial terminal and then launch the cog. The code will then take the value in data var and print on the

## Differences between P1 and P2 Pasm

terminal. Now to the PASM method:

```
23 pub main
24     datavar:= 25           `assign a value to datavar
25
26
27     pst.start(115000)      `start the serial terminal object
28
29     waitcnt(clkfreq*5 +cnt)`hold five sec to open the
30
31     cognew(@asm,@datavar)  ` open a new cog for pasm. where it starts "asm" and
32                             ` the address of the first variable
33     waitcnt(clkfreq+cnt)   ` hold for a second
34
35                             ` print routine
36
37     pst.str(string("answer:"))
38     pst.newline
39     pst.dec(answervar)
40     pst.newline
41
```

Figure 2

The datavar is assigned a value, in this case 256 which is the maximum pasm will handle without extra work. I will tackle that at a later time. We want to keep it simple at this time. This is also because many of the other tutorials I have seen get really complicated very quickly and do not take it in baby steps. I want to make sure that everybody can grasp the concept before getting into complicated code and get lost.

```
42 dat
43
44
45     asm          org      0      `This is the starting point for PASM
46
47     {{ The first item is to move the address of the parameter register "PAR" into
48     a temporary variable and assigne it to the variable in which we will read the in
49     this case the value of datavar in the spin method. }}
50     mov temp_var, par
51     {{ Now we are going to assign the pasm variable, data_var, the address of datavar in
52     the spin method. }}
53     mov data_var, temp_var
54
55     {{ Now we have to move over to the next long to get the address of answervar in the
56     spin object and assign it to answer_var in the pasm code.}}
57     add temp_var, #4
58     {{ Now assign this address to answer_var. }}
59     mov answer_var,temp_var
60
61     {{ Next read the value of datavar (spin object) into the pasm data_var. }}
62     rdlong data_var, par
63
64
65     {{ Finally write it to the answer_var which is spin's answervar for printing. }}
66     wrlong data_var, answer_var
67
68 {{ Reserved variables reserved for PASM's use. }}
69
70     data_var res 1
71     answer_var res 1
72     temp_var res 1
```

Figure 3

The pasm code starts in a “dat” section of spin. The “asm” “org” “0” indicates the beginning of the pasm code. In the cognew there is also an @datavar expression. This tells the pasm code the address of the first variable and that address will be stored in the “par” value. “par” from what I have found means parameter.

There is a very nice webinar done by Jeff Martin in 2009 that explains a lot of information regarding pasm code. I uploaded it to YouTube:  
<https://www.youtube.com/watch?v=OZHuWYW3o1A>.

Starting at:

```
mov temp_var, par
```

This is the mov instruction description:

## MOV

**Instruction:** Set a register to a value.

**MOV** *Destination*,  $\square \# \square$  *Value*

**Result:** *Value* is stored in *Destination*.

*Destination* (d-field) is the register in which to store *Value*.

*Value* (s-field) is a register or a 9-bit literal whose value is stored into *Destination*.

Explanation

MOV copies, or stores, the number in *Value* into *Destination*.

If the WZ effect is specified, the Z flag is set (1) if *Value* equals zero. If the WC effect is specified, the C flag is set to *Value*'s MSB. The result is written to *Destination* unless the NR effect is specified.

So, our first instruction directive will take the address of the spin code datavar variable in the registers and pass it to a temporary variable that we can manipulate. The code is commented so as to follow the progression and I am using full words instead of abbreviations so as one could more easily follow the progression.

```
44
45  asm      org      0      'This is the starting point for PASM
46
47          {{ The first item is to move the address of the parameter register "PAR" into
48          a temporary variable and assigne it to the variable in which we will read the in
49          this case the value of datavar in the spin method. }}
50          mov temp_var, par
51          {{ Now we are going to assign the pasm variable, data_var, the address of datavar in
52          the spin method. }}
53          mov data_var, temp_var
```

Figure 4

Now we have the address of the data\_var which corresponds to datavar in the spin method.

## Differences between P1 and P2 Pasm

```
55      {{ Now we have to move over to the next long to get the address of answervar in the
56      spin object and assign it to answer_var in the pasm code.}}
57      add temp_var, #4
58      {{ Now assign this address to answer_var. }}
59      mov answer_var,temp_var
```

Figure 5

As you can see, we move over and get the address of the spin code answervar variable and assign it's address to the pasm code answer\_var variable. This is done by adding 4 to the temporary variable. Adding 4 moves to the next adjacent long where the answer var is located in the hub.

We are next going to use the rdlong and wrlong directives. The rdlong directive will read from a location and copy the value into a destination field as is shown in the propeller manual listing.

**RDLONG** *Value*, ☐#☐ *Address*

**Result:** Long is stored in *Value*.

*Value* (d-field) is the register to store the long value into.

*Address* (s-field) is a register or a 9-bit literal whose value is the main memory address to read from.

The rdlong goes from right to left. We are reading the value that is in the par register which has the location of datavar and it's contents.

```
60      -      -
61      {{ Next read the value of datavar (spin object) into the pasm data_var. }}
62      rdlong data_var, par
63
```

Figure 6

Lastly, we are going to write the value to the answer\_var location that corresponds with answervar in the spin method and then print the results in a new variable. Note: wrlong works from left to right.

```
65      {{ Finally write it to the answer_var which is spin's answervar for printing. }}
66      wrlong data_var, answer_var
```

Figure 7

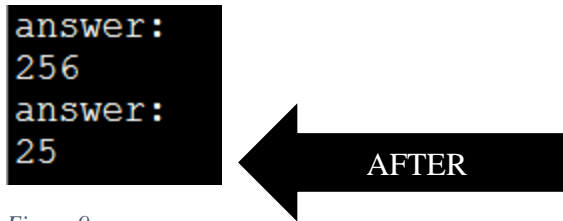
You should get a value on the serial terminal. I used 256 as this is the largest value for a single long, which is four bytes in size.

```
answer:
256
```

Figure 8

Changing the value of datavar to 25 in the spin method to verify.





```

answer:
256
answer:
25

```

AFTER

Figure 9

## RES

**Directive:** Reserve next long(s) for symbol.

⟨*Symbol*⟩ RES ⟨*Count*⟩

- **Symbol** is an optional name for the reserved long in Cog RAM.
- **Count** is the optional number of longs to reserve for *Symbol*. If not specified, RES reserves one long.

RES: We need to reserve space for the pasm variables this is self-explanatory.

Now we can manipulate two variables and print them in succession. This is the new code:

```

1
2
3 {{ Tutorial 2 how to pass two number variables from spin to pasm and back, this works for numbers
4 from 0 to 256, bigger numbers in a later tutorial}}
5
6 CON
7   _clkmode = xtal1 + pll16x
8   _xinfreq = 6_250_000   'MY BOARD AT 100MHZ  DIFFERENT CRYSTAL
9   _xinfreq = 5_000_000   'QUICKSTART 80 MHZ  NORMAL CRYSTAL
10
11 obj
12
13   pst:"parallax serial terminal"
14
15 var
16   long datavar          {{each of these are one long apart. Have to move over one long
17                           so as to access them}}
18   long ansuervar
19   long datavar2
20   long ansuervar2
21

```

Figure 10

```

22 pub main
23   datavar:= 21           'assign a value to datavar
24   datavar2 := 29
25
26   pst.start(115000)      'start the serial terminal object
27
28   waitcnt(clkfreq*5 +cnt)'hold five sec to open the serial terminal
29
30   cognew(@asm,@datavar)  ' open a new cog for pasm. where it starts "asm" and
31                           ' the address of the first variable
32   waitcnt(clkfreq+cnt)   ' hold for a second
33
34                           ' print routine
35
36   pst.str(string("answer:"))
37   pst.newline
38   pst.dec(ansuervar)
39   pst.newline
40   pst.str(string("answer:"))
41   pst.newline
42   pst.dec(ansuervar2)
43   pst.newline
44
45

```

Figure 11

## Differences between P1 and P2 Pasm

```
46
47 dat
48
49 asm      org      0      'This is the starting point for PASM
50
51          {{ The first item is to move the address of the parameter register "PAR" into
52          a temporary variable and assigne it to the variable in which we will read the in
53          this case the value of datavar in the spin method. }}
54          mov temp_var, par
55          {{ Now we are going to assign the pasm variable, data_var, the address of datavar in
56          the spin method. }}
57          mov data_var, temp_var
58          rdlong data_var, temp_var
59          {{ Now we have to move over to the next long to get the address of answervar in the
60          spin object and assign it to answer_var in the pasm code.}}
61          add temp_var, #4
62          {{ Now assign this address to answer_var. }}
63          mov answer_var, temp_var
64          {{write the value to the answervar in spin}}
65          wrlong data_var, answer_var
66          {{go back and get the par address to access the next variable}}
67          mov temp_var, par
```

Figure 12

```
69          {{jump over two longs to get the address of datavar2 in the spin method}}
70          add temp_var, #8
71          {{assign the address}}
72          mov data_var2, temp_var
73          {{read the value}}
74          rdlong data_var2, temp_var
75          {{skip over one long to get answervar in spin}}
76          add temp_var, #4
77          {{assign the address}}
78          mov answer_var2, temp_var
79          {{now write the value to answervar2 in spin}}
80          wrlong data_var2, answer_var2
81
82
83 {{ Reserved variables reserved for PASM's use. }}
84
85 data_var res 1
86 data_var2 res 1
87 answer_var res 1
88 answer_var2 res 1
89 temp_var res 1
```

Figure 13

We have added a couple of items. First a new datavar named datavar2 and a new answervar named answervar2 as well, figure 10 and 14 lines 19 and 20, as their counterparts in the pasm method. In the print area answervar2 has been added also.

```
15 var
16     long datavar      {{each of these are one long apart. Have to move over one long
17                        so as to access them}}
18     long answervar
19     long datavar2
20     long answervar2
21
```

Figure 14

Note the order of the global variables. This will make it easy to find them in the pasm method.

The pasm routine begins just like before and we get the location of datavar from par into the temporary variable and assign the location to data\_var and read the value from par to data\_var.

Now we have to move over a couple of longs to get the new variables and values, figure 13 and 15:



## Differences between P1 and P2 Pasm

```
68      mov temp_var, par
69      {{jump over two longs to get the address of datavar2 in the spin method}}
70      add temp_var, #8
71      {{assign the address}}
72      mov data_var2, temp_var
73      {{read the value}}
74      rdlong data_var2, temp_var
```

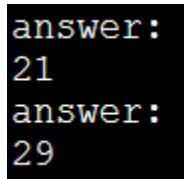
Figure 15

Now we can write the value to the second answer\_var. Remember wrlong is from left to right as opposed to rdlong and other directives which are right to left.

```
76      {{move over one long to get answervar in spin}}
77      add temp_var, #4
78      {{assign the address}}
79      mov answer_var2, temp_var
80      {{now write the value to answervar2 in spin}}
81      wrlong data_var2, answer_var2
```

Figure 16

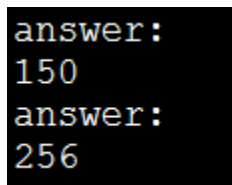
This is what you should see on the serial terminal:



```
answer:
21
answer:
29
```

Figure 17

Changing the two datavar's values:



```
answer:
150
answer:
256
```

Figure 18

It works.

Now that we can get in and out of spin and pasm, I will present some examples of simple math.

I am trying to avoid the jump to really complicated programs with the assumption that the reader has a total comprehension of coding in assembly language of any type. I have found many tutorials do that.

These tutorials were good but confusing when they jump ahead and get very complex. Since I am a teacher, I teach flying and aircraft mechanics, I have to assess the background of each student. Academic learning can be difficult and painful, so if the instructor keeps it simple and explains the concept with easy examples that build up slowly, the student has a better chance of understanding and correlating the subject matter. That results in a much better outcome.

First, we will visit addition.



## Differences between P1 and P2 Pasm

### ADD

**Instruction:** Add two unsigned values.

**ADD** *Value1*, (#) *Value2*

**Result:** Sum of unsigned *Value1* and unsigned *Value2* is stored in *Value1*.

- *Value1* (d-field) is the register containing the value to add to *Value2* and is the destination in which to write the result.
- *Value2* (s-field) is a register or a 9-bit literal whose value is added into *Value1*.

We are going to repeat the above code and make some changes:

Note the global variable name changes, figure 19 lines 10,11 and 12.

```
1 {{basic addition in pasm using the add directive. Page259 propeller manual}}
2 |
3 CON
4 _clkmode = xtall + pll16x
5 _xinfreq = 6_250_000 'MY BOARD AT 100MHZ DIFFERENT CRYSTAL
6 _xinfreq = 5_000_000 'QUICKSTART 80 MHZ NORMAL CRYSTAL
7
8 var
9 'VARIABLE IN THE PAR ADDRESS TO BE PASSED
10 long x
11 long y
12 long product
13 obj
14
15 pst:"parallax serial terminal"
16
17 pub main
18   x := 30
19   y := 45
20 pst.start(115000)
21   waitcnt(clkfreq*5 +cnt) 'hold five sec to open the
22   'serial terminal and enable it
23 cognew(@asm,@x) 'start cog at the first variable address
24   waitcnt(clkfreq*2 +cnt) 'give pasm time to do the work
25
26   pst.str(string("product:"))
27   pst.dec(product*)
28   pst.newline
```

Figure 19

See figure 19 lines 26 and 27.

## Differences between P1 and P2 Pasm

```
29
30 dat
31
32 asm      org
33
34      mov tempvar, par 'get the address of x from par
35      mov xvar, tempvar 'assign the address to the xvar in pasm
36      rdlong xvar, tempvar 'read the value that is in x
37      add tempvar, #4 'move over one long to get y's address
38      mov yvar, tempvar 'assign that address to yvar
39      rdlong yvar, tempvar 'read the value that is in y
40      add tempvar, #4 'move over one long to get the address of product
41      mov productvar, tempvar 'assign the address to productvar
42      add xvar,yvar 'add x and y together answer will be in x
43      wrlong xvar, productvar 'write x into the product variable and print
44
45
46 tempvar long 0
47 xvar long 0
48 yvar long 0
49 productvar long 0
50 flag long 0 |
```

See figure 20 lines 40 to 43. With these changes both variables will be added together.

Figure 20

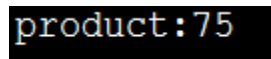


Figure 21

### Subtraction:

Changing line 65 in figure 24 does the trick.

```
1
2
3
4 {{ Tutorial on how to pass a number variable and perform subtraction
5 with the sub directive
6 from spin to pasm and back, this works for numbers
7 from 0 to 256, bigger numbers in a later tutorial}}
8
9
10 CON
11     _clkmode = xtal1 + pll16x
12     _xinfreq = 6_250_000 'MY BOARD AT 100MHZ DIFFERENT CRYSTAL
13     _xinfreq = 5_000_000 'QUICKSTART 80 MHZ NORMAL CRYSTAL
14 obj
15
16     pst:"parallax serial terminal"
17
18 var    'global variables
19     long datavar
20     long ansuervar
21     long subvar
22
```

Figure 22

## Differences between P1 and P2 Pasm

```

23 pub main
24     datavar:= 25          'assign a value to datavar
25     subvar := 10
26
27     pst.start(115000)     'start the serial terminal object
28
29     waitcnt(clkfreq*5 +cnt) 'hold five sec to open the
30
31     cognew(@asm,@datavar) ' open a new cog for pasm. where it starts "asm" and
32                          ' the address of the first variable
33     waitcnt(clkfreq+cnt)  ' hold for a second
34
35                          ' print routine
36
37     pst.str(string("results:"))
38     pst.newline
39     pst.dec(answervar)
40     pst.newline

```

Figure 23

```

41
42 dat
43
44
45     asm          org      0      'This is the starting point for PASM
46
47     {{ The first item is to move the address of the parameter register "PAR" into
48     a temporary variable and assigne it to the variable in which we will read the in
49     this case the value of datavar in the spin method. }}
50     mov temp_var, par
51     {{ Now we are going to assign the pasm variable, data_var, the address of datavar in
52     the spin method. }}
53     mov data_var, temp_var
54
55     {{ Now we have to move over to the next long to get the address of answervar in the
56     spin object and assign it to answer_var in the pasm code.}}
57     add temp_var, #4
58     {{ Now assign this address to answer_var. }}
59     mov answer_var,temp_var
60     add temp_var,#4 'move over to the next long and get the subtraction variable address
61     mov sub_var, temp_var 'assign the address to the variable
62     rdlong sub_var,temp_var 'read the value in that address
63     {{ Next read the value of datavar (spin object) into the pasm data_var. }}
64     rdlong data_var, par 'go back and get the value from the data variable that is in the par register
65     sub data_var,sub_var 'perform the subtraction data-subvar= xxx
66
67     {{ Finally write it to the answer_var which is spin's answervar for printing. }}
68     wrlong data_var, answer_var
69

```

Figure 24

```

70 {{ Reserved variables reserved for PASM's use. }}
71     sub_var res 1
72     data_var res 1
73     answer_var res 1
74     temp_var res 1

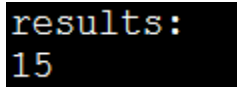
```

Figure 25

What we have done is simply, at lines 60 and 61, added a new variable as well at line 71, these will be the subtraction variables. Next perform the subtraction and then write to our answer variable.

You should get this:

## Differences between P1 and P2 Pasm

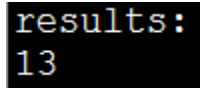


```
results:
15
```

25-10=15

Figure 26

Change subtraction variable to 12.



```
results:
13
```

25-12=13

Figure 27

Multiplication this is from the propeller manual page 380:

```
1 {{Multiplication based on the propeller manual page 380}}
2
3 CON
4 _clkmode = xtall + pll16x
5 _xinfreq = 6_250_000 'MY BOARD AT 100MHZ DIFFERENT CRYSTAL
6 _xinfreq = 5_000_000 'QUICKSTART 80 MHZ NORMAL CRYSTAL
7
8 var
9
10 'VARIABLE IN THE PAR ADDRESS TO BE PASSED
11 long x
12 long y
13 long product
14 obj
15
16 pst:"parallax serial terminal"
17
18 pub main
19     x := 3
20     y := 27
21     pst.start(115000)
22     waitcnt(clkfreq*5 +cnt) 'hold five sec to open the
23     'serial terminal and enable it
24     cognew(@asm,@x) 'start cog at the first variable address |
25     waitcnt(clkfreq*2 +cnt) 'give pasm time to do the work
26
27     pst.str(string("product:"))
28     pst.dec(product~)
29     pst.newline
```

Figure 28

## Differences between P1 and P2 Pasm

```

30
31
32 dat
33 .. Multiply x[15..0] by y[15..0] (y[31..16] must be 0)
34 .. on exit, product in y[31..0]
35
36 asm                org
37
38                mov temp_var, par    'move par to a temporary variable
39                mov x_var, temp_var  'find the x variable
40                rdlong x_var, temp_var 'read in the value from top object
41                add temp_var, #4      'jump to next long which is the address of the
42                'next variable
43                mov y_var, temp_var  'repeat assignment and read in value
44                rdlong y_var, temp_var
45                add temp_var, #4      'jump again to assign the product variable address
46                mov product_var, temp_var

```

Figure 29

```

47
48                multiply shl x_var, #16 'get multiplicand into x[31..16]
49                mov t, #16 'ready for 16 multiplier bits
50                shr y_var, #1 wc 'get initial multiplier bit into c
51 :loop          if_c add y_var, x_var wc 'if c set, add multiplicand to product
52                rcr y_var, #1 wc 'put next multiplier in c, shift prod.
53                djnz t, #:loop 'loop until done
54                wrlong y_var, product_var 'write the product from y[31..0] to the
55                'product variable for the top object
56
57                'multiply_ret ret 'return with product in y[31..0] 'this would be a subroutine
58                'when used in a program
59 temp_var res 1
60 x_var res 1
61 y_var res 1
62 product_var res 1
63 t res 1

```

Figure 30

**product:81** 3\*27=81

Figure 31

Change 27 to 9.

**product:27** 3\*9=27

Figure 32

Basically, we are doing multiplication by addition:

27+27+27=81

3+3+3+3+3+3+3+3+3+3=27

The first operation is to shift left, the multiplicand into x[31..16], line 48.

## SHL

**Instruction:** Shift value left by specified number of bits.

**SHL** *Value*,  $\langle \# \rangle$  *Bits*

---

**Result:** *Value* is shifted left by *Bits*.

- **Value** (d-field) is the register to shift left.
- **Bits** (s-field) is a register or a 5-bit literal whose value is the number of bits to shift left.

Next because this is 16 bit multiplication, so we are going to load a variable with the number 16, line 49: mov t,#16 'ready for 16 multiplier bits.

We are going to shift the carry into y by 1 each time we add the variables. So, on line 50 the first iteration will be loaded. This is done by shifting y right by one to get the carry flag set with the first number that will eventually be the result of the multiplication.

SHR: There is a shift right and shift left these are self-explanatory in the propeller manual as shown. The code will shift left or right by the number specified.

Line 50: shr y\_var,#1 wc 'get initial multiplier bit into c

## SHR

**Instruction:** Shift value right by specified number of bits.

**SHR** *Value*,  $\langle \# \rangle$  *Bits*

---

**Result:** *Value* is shifted right by *Bits*.

- **Value** (d-field) is the register to shift right.
- **Bits** (s-field) is a register or a 5-bit literal whose value is the number of bits to shift right.

Now we are going to ask if the carry flag is set when we add x and y. this will loop until the carry flag is not set and we will loop back and perform the operation again. Each addition will be counted until finished. When completed the carry will be the result of the multiplication. The carry will be discussed in the “if” conditional in the next paragraphs.

Now the loop:

If the carry flag is set, we will loop back and perform an add instruction and check the carry flag after each iteration. This conditional jump will be performed by the DJNZ directive what will evaluate the carry. If the carry in this case is set it will jump back to the beginning of the loop where the RCR instruction will rotate the carry flag, RCR, over into y at the end the value in y will be the answer. Basically, it adds up the carry bits. . If the carry is not set it will NOP, NO OPERATION, and drop out of the loop and go to the next instruction which in this case is to write the results to the variable, product\_var and will be printed.

## Differences between P1 and P2 Pasm

Which in the end of the loop, would be the answer if one did multiplication via the addition process.

RCR:

### RCR

**Instruction:** Rotate C right into value by specified number of bits.

RCR *Value*, <#> *Bits*

---

**Result:** *Value* has *Bits* copies of C rotated right into it.

- **Value** (d-field) is the register in which to rotate C rightwards.
- **Bits** (s-field) is a register or a 5-bit literal whose value is the number of bits of *Value* to rotate C rightwards into.

## CONDITIONAL STATEMENTS:

### IF\_x (Conditions)

Every Propeller Assembly instruction has an optional “condition” field that is used to dynamically determine whether or not it executes when it is reached at run time. The basic syntax for Propeller Assembly instructions is:

<Label> <Condition> *Instruction* *Operands* <Effects>

The optional *Condition* field can contain one of 32 conditions (see Table 3-3) and defaults to IF\_ALWAYS when no condition is specified. The 4-bit **Value** shown for each condition is the value used for the -CON- field in the instruction’s opcode.

This feature, along with proper use of instructions’ optional *Effects* field, makes Propeller Assembly very powerful. Flags can be affected at will and later instructions can be conditionally executed based on the results. Here’s an example:

```
                test    _pins, #$20          wc
                and     _pins, #$38
                shl     t1, _pins
                shr     _pins, #3
                movd    vcfg, _pins
if_nc          mov     dira, t1
if_nc          mov     dirb, #0
if_c           mov     dira, #0
if_c           mov     dirb, t1
```



## Differences between P1 and P2 Pasm

The first instruction, `test _pins, #$20 wc`, performs its operation and adjusts the state of the C flag because the `WC` effect was specified. The next four instructions perform operations that could affect the C flag, but they do not affect it because no `WC` effect was specified. This means that the state of the C flag is preserved since it was last modified by the first instruction. The last four instructions are conditionally executed based on the state of the C flag that was set five instructions prior. Among the last four instructions, the first two `mov` instructions have `if_nc` conditions, causing them to execute only “if not C” (if C = 0). The last two `mov` instructions have `if_c` conditions, causing them to execute only “if C” (if C = 1). In this case, the two pairs of `mov` instructions are executed in a mutually exclusive fashion.

When an instruction’s condition evaluates to `FALSE`, the instruction dynamically becomes a `NOP`, elapsing 4 clock cycles but affecting no flags or registers. This makes the timing of multi-decision code very deterministic.

DJNZ:

### DJNZ

**Instruction:** Decrement value and jump to address if not zero.

DJNZ *Value*, <#> *Address*

**Result:** *Value*-1 is written to *Value*.

- **Value** (d-field) is the register to decrement and test.
- **Address** (s-field) is the register or a 9-bit literal whose value is the address to jump to when the decremented *Value* is not zero.

This directive allows for repetition while decrementing a particular value of choice and when the result is not zero jump to a particular point in the code until the result is zero. At that point the code will drop down to the next instruction in line.

We run the loop until the carry flag is empty. This is repeated addition. Jeff and Dave at Parallax told me that there are many ways to do this. I am working on this myself. Basically, it is repetitive addition and that can be done in a loop until the number of iterations required are completed.

Division:

## Differences between P1 and P2 Pasm

```

1 CON
2 _clkmode = xtall + pll16x
3 _xinfreq = 5_000_000      'QUICKSTART 80 MHZ  NORMAL CRYSTAL
4
5 var
6   long dividend          'VARIABLE IN THE PAR ADDRESS TO BE PASSED
7   long divisor
8   long quotient
9   long remainder
10
11 obj
12   pst : "parallax serial terminal"
13
14 pub main
15   dividend := 211
16   divisor := 6
17   pst.start(115200)
18   waitcnt(clkfreq*5 + cnt) 'hold five sec to open the
19                             'serial terminal and enable it
20   cognew(@asm,@dividend)   'start cog at the first variable address
21   waitcnt(clkfreq + cnt)   'give top object time to catch up to pasm
22
23
24   pst.str(string("quotient:"))
25   pst.dec(quotient)
26   pst.newline
27   pst.str(string("remainder:"))
28   pst.dec(remainder)
29   pst.newline
30

```

Figure 33

```

31
32 dat
33
34
35 asm      org
36
37         mov tempvar, par      'get the par address into the temporary variable
38         rdlong x, tempvar     'read the value into the dividend
39         add tempvar, #4       'move over to the next long to get the divisor variable
40         rdlong y, tempvar     'read the value of the divisor into the variable
41         add tempvar, #4       'move over to the next long to get the quotient address
42
43
44 ' Divide x[31..0] by y[15..0] (y[16] must be 0)
45 ' on exit, quotient is in x[15..0] and remainder is in x[31..16]
46
47 divide   shl y,#15           'get divisor into y[30..15]
48         mov t,#16            'ready for 16 quotient bits
49 :loop    cmpsub x,y          'y <= x? Subtract it, quotient bit in c
50         rcl x,#1            'rotate c into quotient, shift dividend
51         djnz t,#:loop        'loop until done
52
53 ' quotient in x[15..0], ;return if used as a subroutine
54 ' remainder in x[31..16]

```

Figure 34

## Differences between P1 and P2 Pasm

```

55
56      mov    quotientvar,x
57      and    quotientvar,andvar2    'isolate lower 16 bits
58      wrlong quotientvar,tempvar    'write into Spinvar 'quotient'
59
60      mov    remaindervar,x
61      shr    remaindervar, #16    'isolate higher 16 bits
62      add    tempvar,#4    'incr pointer to remainder address
63      wrlong remaindervar,tempvar    'write into Spinvar 'remainder'
64
65 andvar2    long $ffff
66 tempvar    res 1
67 x          res 1
68 y          res 1
69 quotientvar    res 1
70 remaindervar    res 1
71 t          res 1
72

```

Figure 35

```

quotient:35
remainder:1

```

Figure 36

The division will be a continued subtraction algorithm that will subtract the divisor from the dividend until the divisor is either zero or there is a remainder less than the divisor. The answer will now be in the quotient the low bits, with the remainder in the high bits.

On line 47 we are going to shift left the divisor by 15 bits to get it into the high end of y. Then move the number 16 into t because t will be our iterations for the `DNJZ` directive which will perform the loop function 16 iterations. Now the compare and subtract, `cmpsub`, will subtract y from x and see if it is zero, the carry flag will answer the condition. At each iteration we will rotate carry left, `RCL`, by one. At the end of all operations x will have the quotient and y will have the remainder.

## CMPSUB

**Instruction:** Compare two unsigned values and subtract the second if it is lesser or equal.

**CMPSUB** *Value1*, <#> *Value2*

**Result:** Optionally,  $Value1 = Value1 - Value2$ , and Z and C flags = comparison results.

- **Value1** (d-field) is the register containing the value to compare with that of *Value2* and is the destination in which to write the result if a subtraction is performed.
- **Value2** (s-field) is a register or a 9-bit literal whose value is compared with and possibly subtracted from *Value1*.

## RCL

**Instruction:** Rotate C left into value by specified number of bits.

RCL *Value*, ⟨#⟩ *Bits*

**Result:** *Value* has *Bits* copies of C rotated left into it.

- **Value** (d-field) is the register in which to rotate C leftwards.
- **Bits** (s-field) is a register or a 5-bit literal whose value is the number of bits of *Value* to rotate C leftwards into.

The AND operation takes \$FFFF and masks off high bits so as to get the quotient, we later shift the naked remainder by 16 to get the remainder.

## AND – Assembly Language Reference

### AND

**Instruction:** Bitwise AND two values.

AND *Value1*, ⟨#⟩ *Value2*

**Result:** *Value1* AND *Value2* is stored in *Value1*.

- **Value1** (d-field) is the register containing the value to bitwise AND with *Value2* and is the destination in which to write the result.
- **Value2** (s-field) is a register or a 9-bit literal whose value is bitwise ANDed with *Value1*.

Counting up and down:

```

1 {{counting up example, have to slow pasm. Introducing conditionals
2 and jmp command}}
3
4
5
6 CON
7 _clkmode = xtall + pll16x
8 _xinfreq = 6_250_000 'MY BOARD AT 100MHZ
9 _xinfreq = 5_000_000 'QUICKSTART 80 MHZ
10
11 var
12
13     long count
14
15 obj
16
17 pst:"parallax serial terminal"
18
19 pub main
20
21 pst.start(115000)
22     waitcnt(clkfreq*5 +cnt)'hold two sed to open the
23     'serial terminal and enable it
24 cognew(@asm,@count)
25

```

Figure 37

## Differences between P1 and P2 Pasm

```

26 repeat
27   pst.dec(count~)'post clear p 157
28   pst.newline
29   waitcnt(clkfreq +cnt)
30
31 dat
32
33 asm      org
34
35         mov addr, par
36 loop    add value,#1 'counting variable
37 wait    rdlong prev, addr wz 'what is in par??
38         if_nz jmp #wait 'if the value in
39         'par is zero continue to next command
40         'if the value in par "addr" has not been cleared
41         'meaning the value that was put in "value" from
42         'addr which has the address of par "parameter"
43
44         wrlong value, addr
45         'now write the value to the addr which has been assigned
46         'the same address as par and where the address of count in
47         'memory where the spin program can read it then jump back
48         'to the top of the loop and continue after the variable
49         'called count has been cleared to zero
50         jmp #loop
51

```

Figure 38

```

52
53 addr long 0
54 value long 0
55 prev long 0

```

Figure 39

Since spin is much slower than pasm, we have to interrupt pasm so spin can keep up. With that in mind we are going to look at line 27 and 37 to 51.

Line 27:

pst.dec(count~)'post clear p 157.

$Y := X \sim + 2$

The Post-Clear operator in this example clears the variable to 0 (all bits low) after providing its current value for the next operation. In this example if X started out as 5, X~ would provide the current value for the expression (5 + 2) to be evaluated later, then would store 0 in X. The expression 5 + 2 is then evaluated and the result, 7, is stored into Y. After this statement, X equals 0 and Y equals 7.

Since Sign-Extend 7 and Post-Clear are always assignment operators, the rules of Intermediate Assignments apply to them (see page 147).

```

37 wait    rdlong prev, addr wz
38         if_nz jmp #wait

```

Figure 40

## Differences between P1 and P2 Pasm

So, if the line 27 instruction has not cleared, pasm will jump back to the loop until it is cleared then pasm will perform the operation again.

Change line: 36 add to sub and you will have a continuous loop of subtraction.

```
36 loop    add value,#1    'counting variable
```

*Figure 41*

### Arrays:

We are now able to add, subtract, multiply and divide. Basic math skills that we will now take to a next level but in a slow process. Next let's create an array and do some math while learning to populate the array and print selected arrays cells.

### Simple array.

```
1
2 {{basic array populate the array do some simple math }}
3
4
5 CON
6 _clkmode = xtal1 + pll16x
7 _xinfreq = 6_250_000    'MY BOARD AT 100MHZ  DIFFERENT CRYSTAL
8 _xinfreq = 5_000_000    'QUICKSTART 80 MHZ  NORMAL CRYSTAL
9
10 var
11
12
13 long data
14 long array[10] 'global variable array 10 cells long array[0]..array[9]
15
16
17 obj
18
19 pst:"parallax serial terminal"
20
21
22 pub main
23     data := 16
24
25 pst.start(115000)
26     waitcnt(clkfreq*5 +cnt)'hold five sec to open the
27     'serial terminal and enable it
28     cognew(@asm,@data ) 'start cog at the first variable address
29
30
```

*Figure 43*

```
31
32     pst.str(string("array:"))
33     pst.dec(array[1]) 'print the second cell first
34
35     pst.newline
36
37     pst.str(string("array:"))
38     pst.dec(array[0]) 'print the first cell second
39
40     pst.newline
41
42
43
```

*Figure 42*

## Differences between P1 and P2 Pasm

```
44
45
46
47 dat
48
49
50 asm      org 0
51
52      mov tempvar,par    'get the par address into a temporary variable
53
54      mov datvar,tempvar 'assign the address to the datvar in pasm
55
56      rdlong datvar,tempvar 'read in the value of the data variable from spin
57
58      add tempvar,#4      'move over and get the beginning of the array
59
60      mov arrayvar,tempvar 'assign the beginning of the array
61
62      wrlong datvar,arrayvar 'write the value from spin to array[0]
63
64      add arrayvar,#4      'move over to the next array cell
65
66      add datvar, #10      'add 10 to the value in in the data variable from spin
67                          'in this case 16 + 10 = 26
68
69      wrlong datvar,arrayvar 'write the product to the second array cell array[1]
70
```

Figure 44

```
71
72
73
74
75 tempvar long 0
76 datvar long 0
77
78 arrayvar long 10 'global variable array 10 cells long array[0]..array[9]
```

Figure 45

We are going to start as before, and this time have two global variables. One is the data to be passed with a value from spin to pasm. The other is an array that is 10 cells long. That means that each cell will be a long in size.

As you can see in the spin method and the pasm method both are declared. Standard entry to get the addresses and values entered.

The line 56, read in the value to the datavar variable.

Line 62 write it to the first array cell, array[0].

Now to access the second array cell, array[1], we have to move over to the next long, line 64 by adding 4 bytes. Now for a little math to make it interesting we are going to add the littoral number 10 to the variable that is stored in the datavar which is 16. So  $16+10=26$ .

The spin method is going to print them in reverse order which shows that we can manipulate the array.

```
array:26
array:16
```

Figure 46



### P2 PASM SECTION

The P2 does not have an official GUI at this time. Coding is done in “Pnut” or Spin2Gui and a few others. The Pnut is on P2Pasm only. Serial interface has to be coded. Spin2Gui supports .spin2 files, Fastspin, C and Basic. It has a serial terminal as opposed to Pnut.

<https://github.com/totalspectrum/spin2gui/releases>

<https://www.parallax.com/product/64000-es>

The demo code here is in .spin2 and using the Spin2gui

```
1 CON
2   oscmode = $010007f8
3   freq = 160_000_000
4   baud = 230_400
5
6 OBJ
7   ser: "spin/SmartSerial"
8
9 VAR
10  long cog
11  long data
12  long answer
13
14 PUB main
15   clkset(oscmode, freq)
16   ser.start(63, 62, 0, BAUD)
17   pausems(500) ' wait for baud change on host
18
19   ser.str("PASM demo: ")
20
21   data := 12
22   cog := cognew(@startasm, @data)
23
24   ser.str("cog: ")
25   ser.dec(cog)
26
27   '' wait until the COG has finished
28   '' it indicates this by clearing "data"
29   repeat while data <> 0
30
31   '' now print results
32   ser.str(" answer = ")
33   ser.dec(answer)
34   ser.nl
35
36
```

## Differences between P1 and P2 Pasm

```
37 DAT
38 startasm
39     '' assembly program to add 1 to a number
40     '' we start with ptra pointing at the number
41     '' the answer must immediately follow the input data
42     '' in the VAR section
43     rdlong value, ptra
44     add     value, #5
45     add     ptra, #4      ' point to answer
46     wrlong  value, ptra
47     sub     ptra, #4      ' point back to original data
48     wrlong  #0, ptra      ' tell original COG we are done
49
50     '' now shut down
51     cogid   value
52     cogstop value
53
54 value     long    0
55
```

Lines 2,3,4 and 15 are to setup the clock speed and baud rate for the serial terminal. Subsequent lines are similar to regular Spin language. Lines 43 through 48 is where the magic happens.

Instead of using cog ram to share variables we are going to do it with HUB RAM. There is no PAR variable with the P2. There is ptra/ptrb.

We can use the ptra as the start address like in P1 but we are going to work directly with it. So no need to have an intermediary variable to mov things around.

At line 43 we read the value from the ptra directly into the variable that is shared with the .spin2 object and P2 PASM.

So on line 43 we read in the value from spin, then we add 5 to that value , move over to the next long where the answer variable resides and write it to it so it will be printed.

Simple.

## Differences between P1 and P2 Pasm

```
1 CON
2   oscmode = $010007f8
3   freq = 160_000_000
4   baud = 230_400
5
6 OBJ
7   ser: "spin/SmartSerial"
8
9 VAR
10  long cog
11  long data
12  long data2
13  long answer
14  |
15
16 PUB main
17   clkset(oscmode, freq)
18   ser.start(63, 62, 0, BAUD)
19   pausems(500) ' wait for baud change on host
20
21   ser.str(" PASM demo addition: ")
22   ser.nl
23   '   ADD THESE TWO VALUES
24   data := 200
25   data2 := 36500
26   cog := cognew(@startasm, @data)
27
28   ser.str(" cog: ")
29   ser.dec(cog)
30   ser.nl
31
32   ' wait until the COG has finished
33   ' it indicates this by clearing "data"
34   repeat while data <> 0
35
36   ' now print results
37   ser.str(" answer = ")
38   ser.dec(answer)
```

Now here we are going to take two variables and add them, Both will come from the top object:

## Differences between P1 and P2 Pasm

```

33  '' it indicates this by clearing "data"
34  repeat while data <> 0
35
36  '' now print results
37  ser.str(" answer = ")
38  ser.dec(answer)
39  ser.nl
40
41
42  DAT
43  startasm
44      '' assembly program to add numbers
45      '' we start with ptra pointing at the number
46      '' the answer must immediately follow the input data
47      '' in the VAR section
48      rdlong value, ptra      'get first value
49
50      add    ptra, #4         ' move over to pickup next value
51                          'in the next long
52      rdlong value2, ptra     'get second value
53
54      add    value2, value
55      mov    value3,value2    'move to answer var in next long
56      add ptra,#4
57      wrlong value3,ptra      'write it to the pointer
58
59  '*****HERE IS THE QUESTION*****
60      sub    ptra, #8         ' point back to original data????
61      wrlong #0, ptra        ' tell original COG we are done
62
63      '' now shut down
64      cogid  value
65      cogstop value
66
67  value    long    0        'data
68  value2   long    0        'data2
69  value3   long    0        'answer
70

```

Observe line 8 get the first value. Move over 4 long sfor the second. Pick up the second. Add them move over another long to get to the answer to be printed write it and shut down.

## Subtraction”

```

1 CON
2   oscmode = $010007f8
3   freq = 160_000_000
4   baud = 230_400
5
6 OBJ
7   ser: "spin/SmartSerial"
8
9 VAR
10  long cog
11  long data
12  long data2
13  long answer
14
15
16 PUB main
17   clkset(oscmode, freq)
18   ser.start(63, 62, 0, BAUD)
19   pausems(500) ' wait for baud change on host
20
21   ser.str(" PASM demo subtraction: ")
22   ser.nl
23   'SUBTRACT DATA FROM DATA2 if data val is larger it will
24   'generate a negative number
25   data := 35 |
26   data2 := 100
27   cog := cognew(@startasm, @data)
28
29   ser.str(" cog: ")
30   ser.dec(cog)
31   ser.nl

```

```

32
33  '' wait until the COG has finished
34  '' it indicates this by clearing "data"
35  repeat while data <> 0
36
37  '' now print results
38  ser.str(" answer = ")
39  ser.dec(answer)
40  ser.nl
41  ser.nl
42  ' ser.str("answer2 = ")
43  ' ser.dec(answer2)
44  ' ser.nl
45
46
47  DAT
48 startasm
49      '' assembly program to subtract a number
50      '' we start with ptra pointing at the number
51      '' the answer must immediately follow the input data
52      '' in the VAR section
53      rdlong value, ptra      'get first value
54
55      add    ptra, #4        ' move over to pickup next value
56      rdlong value2, ptra    'get second value
57
58      sub    value2, value
59      mov    value3,value2
60      add ptra,#4
61      wrlong value3,ptra
62      sub    ptra, #8        ' point back to original data
63      wrlong #0, ptra        ' tell original COG we are done
64
65      '' now shut down
66      cogid  value
67      cogstop value

```

Multiplication:

## Differences between P1 and P2 Pasm

```
1 CON
2   oscmode = $010007f8
3   freq = 160_000_000
4   baud = 230_400
5
6 OBJ
7   ser: "spin/SmartSerial"
8
9 VAR
10  long cog
11  long data
12  long data2
13  long answer
14
15
16 PUB main
17   clkset(oscmode, freq)
18   ser.start(63, 62, 0, BAUD)
19   pausems(500) ' wait for baud change on host
20
21   ser.str("PASM demo multiplication: ")
22   ser.nl
23   '   MULTIPLY THESE TWO VARIABLES
24   data := 11
25   data2 := 7
26   cog := cognew(@startasm, @data)
27
28   ser.str(" cog: ")
29   ser.dec(cog)
30   ser.nl
31   ' wait until the COG has finished
32   ' it indicates this by clearing "data"
33   repeat while data <> 0
34
35   ' now print results
36   ser.str(" answer = ")
37   ser.dec(answer)
```



## Differences between P1 and P2 Pasm

```
38 ser.nl
39
40
41
42 DAT
43 startasm
44     '' assembly program to multiply numbers
45     '' we start with ptra pointing at the number
46     '' the answer must immediately follow the input data
47     '' in the VAR section
48     rdlong value, ptra      'get first value
49
50     add     ptra, #4        ' move over to pickup next value
51     rdlong  value2, ptra    'get second value
52
53     mul     value2, value
54     mov     value3,value2   'mov results to answer
55     add ptra,#4
56     wrlong  value3,ptra
57     sub     ptra, #8        ' point back to original data
58     wrlong  #0, ptra        ' tell original COG we are done
59
60     '' now shut down
61     cogid   value
62     cogstop value
63
64 value     long    0        'data
65 value2    long    0        'data2
66 value3    long    0        'answer
67
```

Division will be shown on next iteration.

## Differences between P1 and P2 Pasm